

# CHARACTERIZING THE JOIN-IRREDUCIBLE MEDVEDEV DEGREES

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**ABSTRACT.** We characterize the join-irreducible Medvedev degrees as the degrees of complements of Turing ideals, thereby solving a problem posed by Sorbi. We use this characterization to prove that there are Medvedev degrees above the second-least degree that do not bound any join-irreducible degrees above this second-least degree. This solves a problem posed by Sorbi and Terwijn. Finally, we prove that the filter generated by the degrees of closed sets is not prime. This solves a problem posed by Bianchini and Sorbi.

## 1. INTRODUCTION

We present solutions to three problems concerning the Medvedev degrees. A *mass problem* is a set  $\mathcal{A} \subseteq \omega^\omega$ . For mass problems  $\mathcal{A}$  and  $\mathcal{B}$ , we say that  $\mathcal{A}$  *Medvedev reduces* to  $\mathcal{B}$  ( $\mathcal{A} \leq_M \mathcal{B}$ ) if there is a Turing functional  $\Phi$  such that  $\Phi(\mathcal{B}) \subseteq \mathcal{A}$ . That is,  $\Phi(f) \in \mathcal{A}$  for all  $f \in \mathcal{B}$ . We say that  $\mathcal{A}$  and  $\mathcal{B}$  are *Medvedev equivalent* ( $\mathcal{A} \equiv_M \mathcal{B}$ ) if  $\mathcal{A} \leq_M \mathcal{B}$  and  $\mathcal{B} \leq_M \mathcal{A}$ . The equivalence class  $[\mathcal{A}]$  is called the *Medvedev degree* of  $\mathcal{A}$ , and the structure  $\mathfrak{M} = (2^{\omega^\omega} / \equiv_M, \leq_M)$  is called the *Medvedev degrees*. See Sorbi [15] for a good introduction to the theory of the Medvedev degrees.

For  $f, g \in \omega^\omega$ , let  $f \oplus g$  be the function  $(f \oplus g)(2n) = f(n)$  and  $(f \oplus g)(2n+1) = g(n)$ . For  $m \in \omega$  and  $f \in \omega^\omega$ , let  $m \hat{\ } f$  be the function  $(m \hat{\ } f)(0) = m$  and  $(m \hat{\ } f)(n+1) = f(n)$ . In general, ‘ $\hat{\ }$ ’ denotes string concatenation. Functions  $f \in \omega^\omega$  are interpreted as  $\omega$ -length strings when appropriate. For a mass problem  $\mathcal{A}$ , let  $m \hat{\ } \mathcal{A} = \{m \hat{\ } f \mid f \in \mathcal{A}\}$ . Given mass problems  $\mathcal{A}$  and  $\mathcal{B}$ , let  $\mathcal{A} + \mathcal{B} = \{f \oplus g \mid f \in \mathcal{A} \wedge g \in \mathcal{B}\}$  and let  $\mathcal{A} \times \mathcal{B} = 0 \hat{\ } \mathcal{A} \cup 1 \hat{\ } \mathcal{B}$ . Then  $[\mathcal{A}] + [\mathcal{B}] = [\mathcal{A} + \mathcal{B}]$  is the *join* (i.e.,  $\leq_M$ -least upper bound) of  $[\mathcal{A}]$  and  $[\mathcal{B}]$ , while  $[\mathcal{A}] \times [\mathcal{B}] = [\mathcal{A} \times \mathcal{B}]$  is the *meet* (i.e.,  $\leq_M$ -greatest lower bound) of  $[\mathcal{A}]$  and  $[\mathcal{B}]$ . Hence  $\mathfrak{M}$  is a lattice. In fact,  $\mathfrak{M}$  is a distributive lattice, meaning that join and meet distribute over each other:  $\mathbf{a} + (\mathbf{b} \times \mathbf{c}) = (\mathbf{a} + \mathbf{b}) \times (\mathbf{a} + \mathbf{c})$  and  $\mathbf{a} \times (\mathbf{b} + \mathbf{c}) = (\mathbf{a} \times \mathbf{b}) + (\mathbf{a} \times \mathbf{c})$ . Notation for join and meet appears in the literature variously as  $+$ ,  $\times$ , as  $\vee$ ,  $\wedge$ , and confusingly as  $\wedge$ ,  $\vee$ . We choose the  $+$ ,  $\times$  notation to avoid conflict with the logical notation and to match Sorbi and Terwijn [16].

$\mathfrak{M}$  has a least element  $\mathbf{0} = [\omega^\omega]$  (and any  $\mathcal{A}$  containing a recursive function has this degree), a second-least element  $\mathbf{0}' = [\{f \mid f \text{ is Turing jump of } 0\}]$ , and a greatest element  $\mathbf{1} = [\emptyset]$ . (The Medvedev degree  $\mathbf{0}'$  has little to do with  $0'$ , the Turing jump of the 0 function. Here  $\mathbf{0}'$  always refers to the second-least Medvedev degree.)

In any lattice, an element  $\mathbf{a}$  is called *join-reducible* if there are  $\mathbf{x}, \mathbf{y} < \mathbf{a}$  such that  $\mathbf{a} = \mathbf{x} + \mathbf{y}$ . Otherwise  $\mathbf{a}$  is called *join-irreducible*. Dually,  $\mathbf{a}$  is called *meet-reducible* if there are  $\mathbf{x}, \mathbf{y} > \mathbf{a}$  such that  $\mathbf{a} = \mathbf{x} \times \mathbf{y}$ . Otherwise  $\mathbf{a}$  is called *meet-irreducible*. Dymont [3] characterized the meet-reducible Medvedev degrees in the following theorem. Its corollary helps identify meet-irreducible Medvedev degrees.

**Theorem 1.1** ([3]). *A Medvedev degree  $\mathbf{a}$  is meet-reducible if and only if  $\mathbf{a} = [\mathcal{A}]$  for a mass problem  $\mathcal{A}$  for which there are r.e. sets  $V_0, V_1 \subseteq \omega^{<\omega}$  such that*

- $(\forall f \in \mathcal{A})(\exists \sigma \in V_0 \cup V_1)(\sigma \subset f)$ ,
- *The following mass problems are  $\leq_M$ -incomparable:*

$$\{f \in \mathcal{A} \mid (\exists \sigma \in V_0)(\sigma \subset f)\} \text{ and } \{f \in \mathcal{A} \mid (\exists \sigma \in V_1)(\sigma \subset f)\}$$

**Corollary 1.2** ([3]). *If  $\mathcal{A}$  is a mass problem such that  $\sigma \hat{\ } \mathcal{A} \subseteq \mathcal{A}$  for all  $\sigma \in \omega^{<\omega}$ , then  $[\mathcal{A}]$  is meet-irreducible.*

In particular,  $\mathbf{0}'$  is meet-irreducible because  $\sigma \hat{\ } f >_{\mathbf{T}} \mathbf{0}$  whenever  $\sigma \in \omega^{<\omega}$  and  $f >_{\mathbf{T}} \mathbf{0}$ .

The problem of characterizing the join-irreducible Medvedev degrees was posed in [15]. In Section 2, we prove that  $\mathbf{a} \in \mathfrak{M}$  is join-irreducible if and only if  $\mathbf{a} = [\omega^\omega - \mathcal{I}]$  for some Turing ideal  $\mathcal{I}$ .

We have seen that  $\mathfrak{M}$  is a distributive lattice with  $\mathbf{0}$  and  $\mathbf{1}$ . In fact,  $\mathfrak{M}$  is a Brouwer algebra. A *Brouwer algebra* is a distributive lattice with  $\mathbf{0}$  and  $\mathbf{1}$  such that for every  $\mathbf{a}$  and  $\mathbf{b}$  there is a least  $\mathbf{c}$  such that  $\mathbf{a} + \mathbf{c} \geq \mathbf{b}$ . This least  $\mathbf{c}$  is denoted by  $\mathbf{a} \rightarrow \mathbf{b}$ . For mass problems  $\mathcal{A}$  and  $\mathcal{B}$ , define  $\mathcal{A} \rightarrow \mathcal{B} = \{e \hat{\ } g \mid (\forall f \in \mathcal{A})(\Phi_e(f \oplus g) \in \mathcal{B})\}$ . Then  $[\mathcal{A}] \rightarrow [\mathcal{B}] = [\mathcal{A} \rightarrow \mathcal{B}]$ . A Brouwer algebra is dual to a Heyting algebra, but  $\mathfrak{M}$  is proved not to be a Heyting algebra in Sorbi [12].

Brouwer algebras give semantics for propositional logic. For any Brouwer algebra  $\mathfrak{B}$ , a *valuation* is a function  $\nu$ : propositional variables  $\rightarrow \mathfrak{B}$ . A valuation  $\nu$  extends to all propositional formulas  $\varphi$  by defining

$$\begin{aligned} \nu(\varphi \wedge \psi) &= \nu(\varphi) + \nu(\psi), \\ \nu(\varphi \vee \psi) &= \nu(\varphi) \times \nu(\psi), \\ \nu(\varphi \rightarrow \psi) &= \nu(\varphi) \rightarrow \nu(\psi), \text{ and} \\ \nu(\neg \varphi) &= \nu(\varphi) \rightarrow \mathbf{1}. \end{aligned}$$

A propositional formula  $\varphi$  is called *valid* in  $\mathfrak{B}$  if  $\nu(\varphi) = \mathbf{0}$  for every valuation  $\nu$ . Let  $\text{Th}(\mathfrak{B})$  denote the set of propositional formulas valid in  $\mathfrak{B}$ . The axioms of intuitionistic logic are valid in every Brouwer algebra  $\mathfrak{B}$ , so  $\text{IPC} \subseteq \text{Th}(\mathfrak{B}) \subseteq \text{CPC}$  for every Brouwer algebra  $\mathfrak{B}$ . Here IPC denotes intuitionistic logic and CPC denotes classical logic. Logics  $L$  for which  $\text{IPC} \subseteq L \subseteq \text{CPC}$  are called *intermediate logics*.

Providing semantics for propositional logic was one of Medvedev's main motivations behind introducing  $\mathfrak{M}$ , and he proved  $\text{Th}(\mathfrak{M}) = \text{JAN}$  in Medvedev [8]. JAN denotes the logic  $\text{IPC} + \neg p \vee \neg \neg p$  named after Jankov who studied it in Jankov [5]. In any Brouwer algebra  $\mathfrak{B}$ , the quotient of  $\mathfrak{B}$  by the principal filter generated by  $\mathbf{a} \in \mathfrak{B}$  is denoted by  $\mathfrak{B}/\mathbf{a}$ . The quotient  $\mathfrak{B}/\mathbf{a}$  is isomorphic to the interval  $[\mathbf{0}, \mathbf{a}]$  which is a Brouwer algebra under the operations inherited from  $\mathfrak{B}$ . Logics of the form  $\text{Th}(\mathfrak{M}/\mathbf{a})$  have been studied in Skvortsova [10], Sorbi [14], and Sorbi and Terwijn [16]. (Skvortsova and Dymont are the same person. Dymont got married and became Skvortsova.) The results in Section 3 and Section 4 are motivated by the following question which remains open:

**Question 1.3** ([16]). *Is  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$  for all  $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ ?*

Sorbi and Terwijn's study of Question 1.3 in [16] lead them to ask whether every degree  $>_{\mathbf{M}} \mathbf{0}'$  bounds a join-irreducible degree  $>_{\mathbf{M}} \mathbf{0}'$  because a "yes" answer to this question implies a "yes" answer to Question 1.3. However, Sorbi and Terwijn conjectured that there is a degree  $>_{\mathbf{M}} \mathbf{0}'$  that bounds no join-irreducible degree  $>_{\mathbf{M}} \mathbf{0}'$ , and we prove that this is correct in Section 3. In Section 4 we provide slight extensions to some of the results in [14], thereby widening the class of degrees  $\mathbf{a}$  for which  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$  is known.

Lastly, in Section 5 we use techniques similar to those used to characterize the join-irreducible degrees to prove that the filter generated by the degrees of mass problems closed in  $\omega^\omega$  is not prime. This problem was posed in Bianchini and Sorbi [2] and in Sorbi [15].

## 2. CHARACTERIZING THE JOIN-IRREDUCIBLE MEDVEDEV DEGREES

A *Turing ideal* is a set  $\mathcal{I} \subseteq \omega^\omega$  that is closed downward under  $\leq_{\mathbf{T}}$  (i.e.,  $f \in \mathcal{I} \wedge g \leq_{\mathbf{T}} f \rightarrow g \in \mathcal{I}$ ) and closed under  $\oplus$  (i.e.,  $f, g \in \mathcal{I} \rightarrow f \oplus g \in \mathcal{I}$ ). We prove that  $\mathbf{a} \in \mathfrak{M}$  is join-irreducible if and only if  $\mathbf{a} = [\omega^\omega - \mathcal{I}]$  for some Turing ideal  $\mathcal{I}$ . We frequently use the following well-known lemma without mention:

**Lemma 2.1** (see [1] Section III.2). *In a distributive lattice,  $\mathbf{a}$  is join-irreducible if and only if for all  $\mathbf{x}$  and  $\mathbf{y}$ ,  $\mathbf{a} \leq \mathbf{x} + \mathbf{y}$  implies  $\mathbf{a} \leq \mathbf{x}$  or  $\mathbf{a} \leq \mathbf{y}$ . Dually,  $\mathbf{a}$  is meet-irreducible if and only if for all  $\mathbf{x}$  and  $\mathbf{y}$ ,  $\mathbf{a} \geq \mathbf{x} \times \mathbf{y}$  implies  $\mathbf{a} \geq \mathbf{x}$  or  $\mathbf{a} \geq \mathbf{y}$ .*

*Proof.* Suppose  $\mathbf{a}$  is join-irreducible and  $\mathbf{a} \leq \mathbf{x} + \mathbf{y}$ . Then

$$\mathbf{a} = \mathbf{a} \times (\mathbf{x} + \mathbf{y}) = (\mathbf{a} \times \mathbf{x}) + (\mathbf{a} \times \mathbf{y}).$$

Thus  $\mathbf{a} = \mathbf{a} \times \mathbf{x}$  or  $\mathbf{a} = \mathbf{a} \times \mathbf{y}$  which means  $\mathbf{a} \leq \mathbf{x}$  or  $\mathbf{a} \leq \mathbf{y}$ . Conversely, if  $\mathbf{a}$  is join-reducible, then by definition there are  $\mathbf{x}, \mathbf{y} < \mathbf{a}$  with  $\mathbf{a} = \mathbf{x} + \mathbf{y}$ . The proof for the meet-irreducible case is obtained by dualizing the proof for the join-irreducible case.  $\square$

For a mass problem  $\mathcal{A}$ , let  $C(\mathcal{A})$  denote the *Turing upward-closure* of  $\mathcal{A}$ :  $C(\mathcal{A}) = \{f \mid (\exists g \in \mathcal{A})(f \geq_T g)\}$ . A mass problem  $\mathcal{A}$  is called *Turing upward-closed* if  $\mathcal{A} = C(\mathcal{A})$ . The identity functional witnesses  $C(\mathcal{A}) \leq_M \mathcal{A}$  for any mass problem  $\mathcal{A}$ , and if  $\mathcal{A}$  and  $\mathcal{B}$  are mass problems such that  $\mathcal{A}$  is Turing upward-closed, then  $\mathcal{A} \leq_M \mathcal{B}$  if and only if  $\mathcal{B} \subseteq \mathcal{A}$ . Our starting point is the following observation:

**Lemma 2.2** ([15]). *If  $\mathcal{A}$  is a mass problem such that  $[\mathcal{A}]$  is join-irreducible, then  $\omega^\omega - C(\mathcal{A})$  is a Turing ideal.*

*Proof.* We prove the contrapositive. If  $\omega^\omega - C(\mathcal{A})$  is not a Turing ideal, then there are  $f, g \notin C(\mathcal{A})$  with  $f \oplus g \in C(\mathcal{A})$ . This means that  $\{f\}, \{g\} \not\leq_M \mathcal{A}$  but  $\{f\} + \{g\} \geq_M \mathcal{A}$ . Thus  $[\mathcal{A}]$  is join-reducible.  $\square$

The next lemma is the main step in our characterization.

**Lemma 2.3.** *If  $\mathcal{A}$  is a mass problem such that  $[\mathcal{A}]$  is join-irreducible, then  $\mathcal{A} \equiv_M C(\mathcal{A})$*

*Proof.* We prove the contrapositive. Suppose  $\mathcal{A} \not\equiv_M C(\mathcal{A})$ . Then it must be that  $\mathcal{A} \not\leq_M C(\mathcal{A})$ . We find mass problems  $\mathcal{X}$  and  $\mathcal{Y}$  such that  $\mathcal{X}, \mathcal{Y} \not\leq_M \mathcal{A}$  but  $\mathcal{X} + \mathcal{Y} \geq_M \mathcal{A}$ . Thus  $[\mathcal{A}]$  is join-reducible.

To find  $\mathcal{X}$  and  $\mathcal{Y}$ , first find a sequence  $(h_n \mid n \in \omega)$  of functions and a sequence  $(e_n \mid n \in \omega)$  of indices such that

- (i)  $\Phi_{e_n}(h_n) \in \mathcal{A}$  for all  $n \in \omega$ ,
- (ii)  $\Phi_n(h_{2n}) \notin \mathcal{A}$  and  $\Phi_n(h_{2n+1}) \notin \mathcal{A}$  for all  $n \in \omega$ , and
- (iii)  $h_n(0) = \langle n, e_0, e_1, \dots, e_{n-1} \rangle$  for all  $n \in \omega$ .

We find the desired sequences by iterating the following claim:

**Claim.** *If  $\mathcal{A} \not\leq_M C(\mathcal{A})$ , then for every  $e, m \in \omega$  there is an  $h \in C(\mathcal{A})$  such that  $h(0) = m$  and  $\Phi_e(h) \notin \mathcal{A}$ .*

*Proof of claim.* Suppose not. Then there are  $e, m \in \omega$  such that  $h(0) = m$  implies  $\Phi_e(h) \in \mathcal{A}$  for all  $h \in C(\mathcal{A})$ . Thus  $h \mapsto \Phi_e(m \hat{\ } h)$  is a reduction witnessing  $\mathcal{A} \leq_M C(\mathcal{A})$ , a contradiction.  $\square$

Suppose we have  $h_i$  and  $e_i$  for all  $i < n$ . To find  $h_n$  and  $e_n$ , let  $e = \lfloor n/2 \rfloor$  and let  $m = \langle n, e_0, e_1, \dots, e_{n-1} \rangle$ . By the claim, there is an  $h_n \in C(\mathcal{A})$  such that  $h_n(0) = m$  and  $\Phi_e(h_n) \notin \mathcal{A}$ . The fact that  $h_n \in C(\mathcal{A})$  means that there is an  $e_n$  such that  $\Phi_{e_n}(h_n) \in \mathcal{A}$ .

Now set  $\mathcal{X} = \{h_{2n} \mid n \in \omega\}$  and  $\mathcal{Y} = \{h_{2n+1} \mid n \in \omega\}$ . Then  $\Phi_e(\mathcal{X}) \not\subseteq \mathcal{A}$  and  $\Phi_e(\mathcal{Y}) \not\subseteq \mathcal{A}$  for each  $e$  by item (ii). Hence  $\mathcal{X}, \mathcal{Y} \not\leq_M \mathcal{A}$ . The following reduction witnesses  $\mathcal{X} + \mathcal{Y} \geq_M \mathcal{A}$ .

Given  $h$ , decompose  $h$  as  $h = f \oplus g$  and decode  $f(0)$  and  $g(0)$  as  $f(0) = \langle 2n, x_0, x_1, \dots, x_{2n-1} \rangle$  and  $g(0) = \langle 2m+1, y_0, y_1, \dots, y_{2m} \rangle$ . If either  $f(0)$  or  $g(0)$  is not of the required form, then output the 0 function (as such an  $h$  cannot be in  $\mathcal{X} + \mathcal{Y}$ ). Otherwise output  $\Phi_{x_{2m+1}}(g)$  if  $2n > 2m+1$  and output  $\Phi_{y_{2n}}(f)$  if  $2m+1 > 2n$ .

Suppose this reduction is applied to some  $h = h_{2n} \oplus h_{2m+1} \in \mathcal{X} + \mathcal{Y}$ . In this case  $f = h_{2n}$ ,  $g = h_{2m+1}$ , and  $f(0)$  and  $g(0)$  are of the required form by item (iii). So if  $2n > 2m+1$  we output  $\Phi_{e_{2m+1}}(h_{2m+1})$  and if  $2m+1 > 2n$  we output  $\Phi_{e_{2n}}(h_{2n})$ . Both alternatives are in  $\mathcal{A}$  by item (i). Thus  $\mathcal{X} + \mathcal{Y} \geq_M \mathcal{A}$ .  $\square$

**Theorem 2.4.** *A Medvedev degree  $\mathbf{a}$  is join-irreducible if and only if  $\mathbf{a} = [\omega^\omega - \mathcal{I}]$  for some Turing ideal  $\mathcal{I}$ .*

*Proof.* Suppose  $\mathbf{a}$  is join-irreducible, and let  $\mathcal{A}$  be a mass problem such that  $\mathbf{a} = [\mathcal{A}]$ . Then  $\mathcal{I} = \omega^\omega - C(\mathcal{A})$  is a Turing ideal by Lemma 2.2,  $\mathcal{A} \equiv_M C(\mathcal{A})$  by Lemma 2.3, and therefore  $\mathcal{A} \equiv_M C(\mathcal{A}) = \omega^\omega - \mathcal{I}$ . Hence  $\mathbf{a} = [\omega^\omega - \mathcal{I}]$  for the Turing ideal  $\mathcal{I}$ .

Conversely, suppose  $\mathcal{I}$  is a Turing ideal and let  $\mathcal{X}$  and  $\mathcal{Y}$  be mass problems such that  $\mathcal{X}, \mathcal{Y} \not\leq_M \omega^\omega - \mathcal{I}$ . We show that  $\mathcal{X} + \mathcal{Y} \not\leq_M \omega^\omega - \mathcal{I}$ . Observe  $\mathcal{X}, \mathcal{Y} \not\leq \omega^\omega - \mathcal{I}$  for otherwise the identity functional would witness  $\mathcal{X}, \mathcal{Y} \geq_M \omega^\omega - \mathcal{I}$ . Let  $f \in \mathcal{X} \cap \mathcal{I}$  and let  $g \in \mathcal{Y} \cap \mathcal{I}$ , thereby making  $f \oplus g \in (\mathcal{X} + \mathcal{Y}) \cap \mathcal{I}$ . The function  $f \oplus g$  is in  $\mathcal{X} + \mathcal{Y}$ , but it does not compute any member of  $\omega^\omega - \mathcal{I}$ . Therefore  $\mathcal{X} + \mathcal{Y} \not\leq_M \omega^\omega - \mathcal{I}$ . Hence  $[\omega^\omega - \mathcal{I}]$  is join-irreducible.  $\square$

Theorem 2.4 is also valid for the *Muchnik degrees*  $\mathfrak{M}_w$  in place of  $\mathfrak{M}$ , a fact first noticed by Terwijn [17].  $\mathfrak{M}_w$  is defined just as  $\mathfrak{M}$ , but with *Muchnik reducibility* (also called *weak reducibility*)  $\leq_w$  in place of  $\leq_M$ :  $\mathcal{A} \leq_w \mathcal{B}$  if for every  $f \in \mathcal{B}$  there is a  $g \in \mathcal{A}$  such that  $f \geq_T g$ .  $\mathfrak{M}_w$  is a Brouwer algebra with  $+$ ,  $\times$ , and  $\rightarrow$  defined by  $[\mathcal{A}]_w + [\mathcal{B}]_w = [\mathcal{A} + \mathcal{B}]_w$ ,  $[\mathcal{A}]_w \times [\mathcal{B}]_w = [\mathcal{A} \times \mathcal{B}]_w$ , and  $[\mathcal{A}]_w \rightarrow [\mathcal{B}]_w = [\{g \mid (\forall f \in \mathcal{A})(\exists h \in \mathcal{B})(h \leq_T f \oplus g)\}]_w$ . The proof of Lemma 2.2 also works for  $\mathfrak{M}_w$ , and it is easy to check that  $\mathcal{A} \equiv_w C(\mathcal{A})$  for any mass problem  $\mathcal{A}$  (i.e., the  $\mathfrak{M}_w$  analogue of Lemma 2.3 is trivial). This gives the forward direction of Theorem 2.4 for  $\mathfrak{M}_w$ . The proof of the reverse direction of Theorem 2.4 also works for  $\mathfrak{M}_w$ .

### 3. DEGREES THAT BOUND NO JOIN-IRREDUCIBLE DEGREES $>_M \mathbf{0}'$

Recall that JAN is the intermediate logic  $\text{IPC} + \neg p \vee \neg \neg p$ . The results of this section and the next are motivated by Question 1.3: is  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$  for every  $\mathbf{a} >_M \mathbf{0}'$ ?

$\text{Th}(\mathfrak{M}/\mathbf{0}') = \text{CPC}$  because  $\mathfrak{M}/\mathbf{0}' \cong [\mathbf{0}, \mathbf{0}'] = \{\mathbf{0}, \mathbf{0}'\}$ . In fact,  $\mathbf{0}'$  is the only degree for which  $\text{Th}(\mathfrak{M}/\mathbf{0}') = \text{CPC}$ . This is because if  $\mathbf{a} >_M \mathbf{0}'$ , then  $\mathbf{0}' \rightarrow \mathbf{a} = \mathbf{a}$ , hence  $\mathbf{0}' \times (\mathbf{0}' \rightarrow \mathbf{a}) = \mathbf{0}'$ . Thus let  $p = \mathbf{0}'$  to see that the formula  $p \vee \neg p$  is not valid in  $\text{Th}(\mathfrak{M}/\mathbf{a})$ .

Furthermore, if  $\mathbf{a} >_M \mathbf{0}'$ , then we cannot have  $\text{Th}(\mathfrak{M}/\mathbf{a}) \supseteq \text{JAN}$ . It is an easy check that in any Brouwer algebra  $\mathfrak{B}$  with meet-irreducible  $\mathbf{0}$  (such as the algebras  $\mathfrak{M}/\mathbf{a}$ ),  $\neg p \vee \neg \neg p \in \text{Th}(\mathfrak{B})$  if and only if  $\mathbf{1}$  is join-irreducible. However, if  $\mathbf{a} >_M \mathbf{0}'$  is join-irreducible, then  $\text{Th}(\mathfrak{M}/\mathbf{a}) = \text{JAN}$  [14]. Thus if  $\mathbf{a} >_M \mathbf{0}'$  and  $\text{Th}(\mathfrak{M}/\mathbf{a}) \supseteq \text{JAN}$ , then  $\neg p \vee \neg \neg p \in \text{Th}(\mathfrak{M}/\mathbf{a})$  which implies that  $\mathbf{a}$  is join-irreducible which implies that  $\text{Th}(\mathfrak{M}/\mathbf{a}) = \text{JAN}$ . Thus a “no” answer to Question 1.3 must yield a degree  $\mathbf{a}$  such that  $\text{Th}(\mathfrak{M}/\mathbf{a}) \not\subseteq \text{JAN}$  and  $\text{JAN} \not\subseteq \text{Th}(\mathfrak{M}/\mathbf{a})$ .

The following theorem shows that to give a “yes” answer to Question 1.3 it suffices to show that every  $\mathbf{a} >_M \mathbf{0}'$  bounds a finite meet of join-irreducible degrees  $>_M \mathbf{0}'$ .

**Theorem 3.1** ([14]). *If  $\mathbf{a}$  is a degree such that  $\mathbf{a} \geq_M \prod_{i=0}^n \mathbf{d}_i$  for join-irreducible degrees  $\mathbf{d}_i >_M \mathbf{0}'$ ,  $i \leq n$ , then  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$ .*

(The above theorem is stated more generally in [14]. Each degree  $\mathbf{d}_i$  for  $i \leq n$  is allowed to be either join-irreducible or  $\mathfrak{D}\mathfrak{e}$ -irreducible. See the parenthetical discussion following Theorem 4.1 for the definition of  $\mathfrak{D}\mathfrak{e}$ -irreducible and an explanation of why we do not consider such degrees here. Theorem 4.1 is a restatement of [14] Theorem 2.11 which is the main tool used to prove Theorem 3.1.)

The degrees of the mass problems  $\mathcal{B}_f = \{g \mid g \not\leq_T f\}$  play an important role in the study of Question 1.3. The following lemma from Sorbi [13] encapsulates the properties of the  $[\mathcal{B}_f]$ 's that we need in this section and the next.

**Lemma 3.2** ([13]).

- (i) *Every  $[\mathcal{B}_f]$  is join-irreducible.*
- (ii) *Every  $\sum_{i=1}^n [\mathcal{B}_{f_i}]$  is meet-irreducible.*

- (iii) Let  $V$  and  $J$  be finite sets and let  $U_v$  and  $I_j$  be finite sets for each  $v \in V$  and  $j \in J$ . Let  $\mathbf{x}_u^v$  and  $\mathbf{y}_i^j$  be degrees of the form  $[\mathcal{B}_f]$  for every  $v \in V$ ,  $u \in U_v$ ,  $j \in J$ , and  $i \in I_j$ . Let  $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$  and  $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$ . Then  $\mathbf{a} \leq_M \mathbf{b}$  if and only if

$$(\forall v \in V)(\exists j \in J)(\forall i \in I_j)(\exists u \in U_v)(\mathbf{x}_u^v \leq_M \mathbf{y}_i^j).$$

- (iv) In the notation of item (iii),

$$\mathbf{a} \rightarrow \mathbf{b} = \sum \left\{ \prod_{i \in I_j} \mathbf{y}_i^j \mid (\forall v \in V) \left( \prod_{i \in I_j} \mathbf{y}_i^j \not\leq_M \prod_{u \in U_v} \mathbf{x}_u^v \right) \right\}$$

(where the empty join is  $\mathbf{0}$ ).

*Proof.* Item (i) is by Theorem 2.4 and item (ii) is by Corollary 1.2. Item (iv) is proved in [13]. Item (iii) follows from item (iv) because  $\mathbf{a} \leq_M \mathbf{b}$  if and only if  $\mathbf{b} \rightarrow \mathbf{a} = \mathbf{0}$ .  $\square$

In [16] it is asked if every degree  $\mathbf{a} >_M \mathbf{0}'$  bounds a join-irreducible degree  $>_M \mathbf{0}'$ , and it is conjectured that this is not the case based on the evidence provided by the following theorem.

**Theorem 3.3** ([16]). *There is a degree  $\mathbf{a} >_M \mathbf{0}'$  such that  $\mathbf{a} \not\leq_M [\mathcal{B}_f]$  for every  $f >_T \mathbf{0}$ .*

Our characterization of the join-irreducible degrees implies that every join-irreducible degree  $>_M \mathbf{0}'$  bounds some degree  $[\mathcal{B}_f]$  with  $f >_T \mathbf{0}$ . Thus the conjecture is correct.

**Corollary 3.4** (to Theorem 2.4). *If  $\mathbf{a} >_M \mathbf{0}'$  is join-irreducible, then  $\mathbf{a} \geq_M [\mathcal{B}_f]$  for some  $f >_T \mathbf{0}$ .*

*Proof.* If  $\mathbf{a}$  is join-irreducible, then, by Theorem 2.4,  $\mathbf{a} = [\omega^\omega - \mathcal{I}]$  for some Turing ideal  $\mathcal{I}$ . If  $[\omega^\omega - \mathcal{I}] >_M \mathbf{0}'$ , then  $\mathcal{I}$  contains some function  $f >_T \mathbf{0}$ . Thus  $\omega^\omega - \mathcal{I} \subseteq \mathcal{B}_f$ . Hence  $\mathbf{a} = [\omega^\omega - \mathcal{I}] \geq_M [\mathcal{B}_f]$ .  $\square$

**Theorem 3.5.** *There is a degree  $\mathbf{a} >_M \mathbf{0}'$  such that every degree  $\mathbf{x}$  with  $\mathbf{0}' <_M \mathbf{x} \leq_M \mathbf{a}$  is join-irreducible.*

*Proof.* By Theorem 3.3, let  $\mathbf{a} >_M \mathbf{0}'$  be such that  $\mathbf{a} \not\leq_M [\mathcal{B}_f]$  for every  $f >_T \mathbf{0}$ . This  $\mathbf{a}$  is the desired degree because, by Corollary 3.4, if  $\mathbf{a} \geq_M \mathbf{x}$  for some join-irreducible  $\mathbf{x} >_M \mathbf{0}'$ , then  $\mathbf{a} \geq_M [\mathcal{B}_f]$  for some  $f >_T \mathbf{0}$ .  $\square$

The degree  $\mathbf{a}$  satisfying Theorem 3.3 was constructed by diagonalization in [16]. We can give somewhat more concrete examples of degrees satisfying Theorem 3.3 and Theorem 3.5. Recall the following definitions. Functions  $f, g >_T \mathbf{0}$  are a *Turing minimal pair* if, for all  $h$ ,  $h \leq_T f, g$  implies  $h \leq_T \mathbf{0}$ . A function  $f$  has *minimal Turing degree* if, for all  $h$ ,  $h <_T f$  implies  $h \leq_T \mathbf{0}$ . Minimal pairs and minimal degrees exist. In fact, there are continuum many distinct minimal Turing degrees. See Lerman [6] Section II.4 and Section V.2.

**Theorem 3.6.** *If  $f$  and  $g$  are a minimal pair, then the degree  $\mathbf{a} = [\mathcal{B}_f] \times [\mathcal{B}_g]$  witnesses Theorem 3.5.*

*Proof.* Let  $f$  and  $g$  be a minimal pair. Then  $[\mathcal{B}_f], [\mathcal{B}_g] >_M \mathbf{0}'$  because  $f, g >_T \mathbf{0}$ . Thus  $[\mathcal{B}_f] \times [\mathcal{B}_g] >_M \mathbf{0}'$  because  $\mathbf{0}'$  is meet-irreducible by Corollary 1.2. To show that  $[\mathcal{B}_f] \times [\mathcal{B}_g]$  bounds no join-irreducible degree  $>_M \mathbf{0}'$ , it suffices by Corollary 3.4 to show that  $[\mathcal{B}_f] \times [\mathcal{B}_g]$  bounds no  $[\mathcal{B}_h]$  for  $h >_T \mathbf{0}$ . This is true because  $f, g$  is a minimal pair, so for any  $h >_T \mathbf{0}$ , either  $h \not\leq_T f$  or  $h \not\leq_T g$ . Thus either  $h \in \mathcal{B}_f$  or  $h \in \mathcal{B}_g$  which means  $\mathcal{B}_f \times \mathcal{B}_g$  contains a function  $\equiv_T h$ .  $\mathcal{B}_h$  contains no function  $\leq_T h$ , therefore  $\mathcal{B}_f \times \mathcal{B}_g \not\leq_M \mathcal{B}_h$ .  $\square$

We can extend the idea behind Theorem 3.6 to find a degree  $\mathbf{a} >_M \mathbf{0}'$  that does not bound any finite meet of join-irreducible degrees  $>_M \mathbf{0}'$ . Several of our examples in this section and the next are of the form  $[\bigcup_{i \in \omega} i \hat{\ } \mathcal{D}_i]$  for mass problems  $\mathcal{D}_i$ ,  $i \in \omega$ .

**Lemma 3.7.** *Let  $\mathbf{d} = [\bigcup_{i \in \omega} i \hat{\ } \mathcal{D}_i]$  where  $[\mathcal{D}_i] >_M \mathbf{0}'$  for each  $i \in \omega$ . Then  $\mathbf{d} >_M \mathbf{0}'$ .*

*Proof.* Suppose for a contradiction that  $\Phi$  is a reduction witnessing  $\mathbf{d} \leq_M \mathbf{0}'$  (i.e.,  $\Phi(f) \in \bigcup_{i \in \omega} i \hat{\ } \mathcal{D}_i$  for all  $f \succ_T 0$ ). Let  $\sigma \in \omega^{<\omega}$  be such that  $\Phi(\sigma)(0) \downarrow$  and let  $i = \Phi(\sigma)(0)$ . Then  $f \mapsto \Phi(\sigma \hat{\ } f)$  is a reduction witnessing  $\mathbf{0}' \geq_M [\mathcal{D}_i]$ , contradicting  $[\mathcal{D}_i] \succ_M \mathbf{0}'$ .  $\square$

**Theorem 3.8.** *There is a degree  $\mathbf{a} \succ_M \mathbf{0}'$  such that no degree  $\mathbf{x}$  with  $\mathbf{0}' <_M \mathbf{x} \leq_M \mathbf{a}$  is of the form  $\prod_{i=0}^n \mathbf{d}_i$  for join-irreducible degrees  $\mathbf{d}_i \succ_M \mathbf{0}'$ ,  $i \leq n$ .*

*Proof.* By Corollary 3.4, it suffices to find a degree  $\mathbf{a} \succ_M \mathbf{0}'$  which is not above any degree of the form  $\prod_{i=0}^n [\mathcal{B}_{f_i}]$  where  $f_i \succ_T 0$  for each  $i \leq n$ . Let  $\{g_i \mid i \in \omega\}$  be a countable collection of functions all of distinct minimal Turing degree. Let  $\mathcal{A} = \bigcup_{i \in \omega} i \hat{\ } \mathcal{B}_{g_i}$  and put  $\mathbf{a} = [\mathcal{A}]$ . Lemma 3.7 proves that  $\mathbf{a} \succ_M \mathbf{0}'$ .

Now consider any degree  $\prod_{i=0}^n [\mathcal{B}_{f_i}]$ , where  $f_i \succ_T 0$  for each  $i \leq n$ . There is a  $j \in \omega$  such that  $g_j \not\leq_T f_i$  for each  $i \leq n$ . Thus for this  $j$ ,  $[\mathcal{B}_{g_j}] \not\leq_M [\mathcal{B}_{f_i}]$  for each  $i \leq n$ . Therefore  $[\mathcal{B}_{g_j}] \not\leq_M \prod_{i=0}^n [\mathcal{B}_{f_i}]$  because  $[\mathcal{B}_{g_j}]$  is meet-irreducible. Clearly  $[\mathcal{B}_{g_j}] \geq_M \mathbf{a}$ , so  $\mathbf{a} \not\leq_M \prod_{i=0}^n [\mathcal{B}_{f_i}]$  as well.  $\square$

For mass problems  $\mathcal{A}_i$ ,  $i \in \omega$ , the Medvedev degree  $[\bigcup_{i \in \omega} i \hat{\ } \mathcal{A}_i]$  is not in general the greatest lower bound of the degrees  $[\mathcal{A}_i]$ ,  $i \in \omega$ . Such greatest lower bounds need not even exist. For example, the degrees  $[\mathcal{B}_{g_i}]$ ,  $i \in \omega$  from Theorem 3.8 do not have a greatest lower bound. This follows from results in Dymont [4] which studies when countable collections of degrees have least upper bounds and greatest lower bounds.

If  $\mathbf{a}$  is a degree such that  $\mathbf{a} \not\leq_M \mathbf{d}$  for all join-irreducible  $\mathbf{d} \succ_M \mathbf{0}'$ , then  $\mathbf{a} \rightarrow \mathbf{d} = \mathbf{d}$  for all join-irreducible  $\mathbf{d} \succ_M \mathbf{0}'$ . The degree  $\mathbf{a}$  constructed in Theorem 3.8 enjoys a similar property.

**Theorem 3.9.** *There is a degree  $\mathbf{a} \succ_M \mathbf{0}'$  such that  $\mathbf{a} \rightarrow \prod_{i=0}^n \mathbf{d}_i = \prod_{i=0}^n \mathbf{d}_i$  whenever  $\mathbf{d}_i \succ_M \mathbf{0}'$  and is join-irreducible for each  $i \leq n$ .*

*Proof.* As in Theorem 3.8, let  $\{g_i \mid i \in \omega\}$  be a countable collection of functions all of distinct minimal Turing degree, let  $\mathcal{A} = \bigcup_{i \in \omega} i \hat{\ } \mathcal{B}_{g_i}$ , and put  $\mathbf{a} = [\mathcal{A}]$ . Suppose  $\mathbf{d}_i \succ_M \mathbf{0}'$  and is join-irreducible for each  $i \leq n$ . By Theorem 2.4, for each  $i \leq n$  let  $\mathcal{I}_i \subseteq \omega^\omega$  be a Turing ideal such that  $\mathbf{d}_i = [\omega^\omega - \mathcal{I}_i]$ . Thus  $\prod_{i=0}^n \mathbf{d}_i = [\bigcup_{i=0}^n i \hat{\ } (\omega^\omega - \mathcal{I}_i)]$  and

$$\mathbf{a} \rightarrow \prod_{i=0}^n \mathbf{d}_i = \left[ \left\{ e \hat{\ } g \mid \left( \forall f \in \mathcal{A} \right) \left( \Phi_e(f \oplus g) \in \bigcup_{i=0}^n i \hat{\ } (\omega^\omega - \mathcal{I}_i) \right) \right\} \right].$$

We now describe a reduction witnessing  $\mathbf{a} \rightarrow \prod_{i=0}^n \mathbf{d}_i \geq_M \prod_{i=0}^n \mathbf{d}_i$ .

Given  $e \hat{\ } g$ , for each  $i \leq n+1$  search for a string  $i \hat{\ } \sigma_i$  such that  $\Phi_e((i \hat{\ } \sigma_i) \oplus g)(0) \downarrow$ . If there is a  $k \leq n$  such that

$$\Phi_e((i \hat{\ } \sigma_i) \oplus g)(0) = \Phi_e((j \hat{\ } \sigma_j) \oplus g)(0) = k$$

for two distinct  $i, j \leq n+1$ , choose the least such  $k$  and output  $k \hat{\ } g$ . Otherwise output 0.

Suppose we apply this reduction to  $e \hat{\ } g \in \mathcal{A} \rightarrow \bigcup_{i=0}^n i \hat{\ } (\omega^\omega - \mathcal{I}_i)$ .  $\Phi_e(f \oplus g)$  must be total for each  $f \in \mathcal{A}$ , and for each  $i \in \omega$  there is an  $f \in \mathcal{A}$  with  $f(0) = i$ . Thus for each  $i \leq n+1$  the search finds a string  $i \hat{\ } \sigma_i$  such that  $\Phi_e((i \hat{\ } \sigma_i) \oplus g)(0) \downarrow$ . Moreover, each  $i \hat{\ } \sigma_i$  can be extended to a function in  $\mathcal{A}$ , so  $\Phi_e((i \hat{\ } \sigma_i) \oplus g)(0) \leq n$  for each  $i \leq n+1$ . Therefore there is a least  $k \leq n$  for which there are distinct  $i, j \leq n+1$  with  $\Phi_e((i \hat{\ } \sigma_i) \oplus g)(0) = \Phi_e((j \hat{\ } \sigma_j) \oplus g)(0) = k$ . The reduction outputs  $k \hat{\ } g$ , so we must show that  $k \hat{\ } g \in \bigcup_{i=0}^n i \hat{\ } (\omega^\omega - \mathcal{I}_i)$  which means we must show that  $g \notin \mathcal{I}_k$ . Suppose for a contradiction that  $g \in \mathcal{I}_k$ . The functions  $g_i$  and  $g_j$  have distinct minimal degree, so either  $g \not\leq_T g_i$  or  $g \not\leq_T g_j$  ( $g \succ_T 0$  because  $\mathbf{a} \not\leq_M \prod_{i=0}^n \mathbf{d}_i$  by Theorem 3.8). For the sake of argument, suppose  $g \not\leq_T g_i$ . Then  $\sigma_i \hat{\ } g \not\leq_T g_i$  as well, so  $\sigma_i \hat{\ } g \in \mathcal{B}_{g_i}$  and  $i \hat{\ } \sigma_i \hat{\ } g \in \mathcal{A}$ . However,  $\Phi_e((i \hat{\ } \sigma_i \hat{\ } g) \oplus g) \in k \hat{\ } (\omega^\omega - \mathcal{I}_k)$  by the choice of  $i \hat{\ } \sigma_i$ . This cannot be because  $(i \hat{\ } \sigma_i \hat{\ } g) \oplus g \in \mathcal{I}_k$ , thus anything it computes is also in  $\mathcal{I}_k$ .  $\square$

By Corollary 4.6 below, the degree  $\mathbf{a} = [\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{B}_{g_i}]$  used to witness Theorem 3.8 and Theorem 3.9 satisfies  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$  and so does any degree that bounds it. There are, however, degrees  $>_{\mathbf{M}} \mathbf{0}'$  that do not bound any degree of the form  $[\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i]$  where  $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$  and is join-irreducible for each  $i \in \omega$ .

**Theorem 3.10.** *There is a degree  $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$  such that  $\mathbf{a} \not\leq_{\mathbf{M}} [\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i]$  whenever  $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$  and is join-irreducible for each  $i \in \omega$ .*

*Proof.* Let  $\mathcal{D}_i$  be such that  $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$  and is join-irreducible for each  $i \in \omega$ . By Corollary 3.4, for every  $i \in \omega$  there is an  $f_i >_{\mathbf{T}} 0$  such that  $\mathcal{D}_i \geq_{\mathbf{M}} \mathcal{B}_{f_i}$ . The mass problem  $\mathcal{B}_{f_i}$  is Turing upward-closed for each  $i \in \omega$ , so  $\mathcal{D}_i \subseteq \mathcal{B}_{f_i}$  for each  $i \in \omega$ . Thus  $\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i \subseteq \bigcup_{i \in \omega} i \hat{\wedge} \mathcal{B}_{f_i}$ . Hence it suffices to find a degree  $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$  that does not bound any degree of the form  $[\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{B}_{f_i}]$ , where  $f_i >_{\mathbf{T}} 0$  for each  $i \in \omega$ .

We use the same construction used in [16] to prove Theorem 3.3. Build mass problems  $\mathcal{A}_s \subseteq \{g \mid g >_{\mathbf{T}} 0\}$  such that  $\{g \mid g >_{\mathbf{T}} 0\} - \mathcal{A}_s$  is finite for each  $s \in \omega$ . Set  $\mathcal{A}_0 = \{g \mid g >_{\mathbf{T}} 0\}$ . At stage  $s + 1$ , choose  $h_s >_{\mathbf{T}} 0$  such that  $h_s$  does not compute any of the (finitely many) functions in  $\{g \mid g >_{\mathbf{T}} 0\} - \mathcal{A}_s$ . If  $\Phi_s(h_s)$  is total and  $>_{\mathbf{T}} 0$ , let  $g_s = \Phi_s(h_s)$  and set  $\mathcal{A}_{s+1} = \mathcal{A}_s - \{g_s\}$ . Otherwise set  $\mathcal{A}_{s+1} = \mathcal{A}_s$ . Put  $\mathcal{A} = \bigcap_{s \in \omega} \mathcal{A}_s$  and put  $\mathbf{a} = [\mathcal{A}]$ .

To see  $\mathbf{a} >_{\mathbf{M}} \mathbf{0}'$ , observe that by construction  $\Phi_s(h_s) \notin \mathcal{A}$  for each  $s \in \omega$ . Now let  $f_i >_{\mathbf{T}} 0$  for each  $i \in \omega$ . We need to show that  $\Phi_e(\mathcal{A}) \not\subseteq \bigcup_{i \in \omega} i \hat{\wedge} \mathcal{B}_{f_i}$  for every index  $e$ . To do this, we first show that the functions in  $\{g \mid g >_{\mathbf{T}} 0\} - \mathcal{A}$  have distinct Turing degree. Suppose that  $g_i$  leaves  $\mathcal{A}$  at stage  $i + 1$  and  $g_j$  leaves  $\mathcal{A}$  at stage  $j + 1$  for  $i + 1 < j + 1$  (i.e., at stage  $i + 1$  we had  $\Phi_i(h_i) = g_i >_{\mathbf{T}} 0$ , and at stage  $j + 1$  we had  $\Phi_j(h_j) = g_j >_{\mathbf{T}} 0$ ). Then  $g_i \not\leq_{\mathbf{T}} g_j$  because otherwise  $g_i \leq_{\mathbf{T}} g_j \leq_{\mathbf{T}} h_j$ , contradicting that  $h_j$  was chosen  $\not\leq_{\mathbf{T}} g_i$  at stage  $j + 1$ . Now suppose  $\Phi_e(g)$  is total for all  $g \in \mathcal{A}$ . Fix any  $\sigma \in \omega^{<\omega}$  such that  $\Phi_e(\sigma)(0) \downarrow$ , and let  $n$  be such that  $\Phi_e(\sigma)(0) = n$ .  $\mathcal{A}$  is missing at most one function  $\equiv_{\mathbf{T}} f_n$ , so let  $g \in \mathcal{A}$  be such that  $\sigma \subset g$  and  $g \equiv_{\mathbf{T}} f_n$ . Then  $\Phi_e(g)(0) = n$ , but  $\Phi_e(g) \notin n \hat{\wedge} \mathcal{B}_{f_n}$ . Hence  $\Phi_e(\mathcal{A}) \not\subseteq \bigcup_{i \in \omega} i \hat{\wedge} \mathcal{B}_{f_i}$ .  $\square$

**Question 3.11.** Let  $\mathbf{a}$  be the degree constructed in Theorem 3.10. Does  $\mathbf{a} \rightarrow [\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i] = [\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i]$  whenever  $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$  and is join-irreducible for each  $i \in \omega$ ? Is  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$ ?

Finally, we note that the answer to Question 1.3 is “no” for  $\mathfrak{M}_w$  in place of  $\mathfrak{M}$ . Let  $f >_{\mathbf{T}} 0$  have minimal Turing degree, and let  $\mathbf{a} = [\mathcal{B}_f]_w$ . Then, in  $\mathfrak{M}_w$ ,  $[\mathbf{0}, \mathbf{a}] = \{\mathbf{0}, \mathbf{0}', \mathbf{a}\}$  and  $\text{JAN} \not\subseteq \text{Th}(\mathfrak{M}_w/\mathbf{a}) \not\subseteq \text{CPC}$ .

#### 4. NEW DEGREES WHOSE CORRESPONDING LOGIC IS CONTAINED IN JAN

We extend Theorem 3.1 by proving  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$  for degrees  $\mathbf{a}$  such that  $\mathbf{a} \geq_{\mathbf{M}} [\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i]$  for some collection of join-irreducible degrees  $[\mathcal{D}_i] >_{\mathbf{M}} \mathbf{0}'$ ,  $i \in \omega$ .

A propositional formula is called *positive* if the connective ‘ $\neg$ ’ does not appear in it. For a logic  $L$  let  $L^+$  denote the positive formulas in  $L$ , and for a Brouwer algebra  $\mathfrak{B}$  let  $\text{Th}^+(\mathfrak{B})$  denote the set of positive formulas valid in  $\mathfrak{B}$ . JAN is the maximum intermediate logic  $L$  for which  $L^+ = \text{IPC}^+$  [5]. This means that  $L^+ = \text{IPC}^+$  implies  $L \subseteq \text{JAN}$  for any intermediate logic  $L$ . Thus  $\text{Th}^+(\mathfrak{B}) = \text{IPC}^+$  implies  $\text{Th}(\mathfrak{B}) \subseteq \text{JAN}$  for any Brouwer algebra  $\mathfrak{B}$ .

Let  $\mathfrak{B}_1$  and  $\mathfrak{B}_2$  be Brouwer algebras. An injection  $f: \mathfrak{B}_1 \rightarrow \mathfrak{B}_2$  is called a *B-embedding* if it preserves  $\mathbf{0}$ ,  $\mathbf{1}$ ,  $+$ ,  $\times$ , and  $\rightarrow$  (and therefore also  $\neg$ ). An injection  $f: \mathfrak{B}_1 \rightarrow \mathfrak{B}_2$  is called a *B<sup>+</sup>-embedding* if it preserves  $\mathbf{0}$ ,  $+$ ,  $\times$ , and  $\rightarrow$  (but not necessarily  $\mathbf{1}$  or  $\neg$ ). If  $\mathfrak{B}_1$  B-embeds into  $\mathfrak{B}_2$ , then  $\text{Th}(\mathfrak{B}_2) \subseteq \text{Th}(\mathfrak{B}_1)$ , and if  $\mathfrak{B}_1$  B<sup>+</sup>-embeds into  $\mathfrak{B}_2$ , then  $\text{Th}^+(\mathfrak{B}_2) \subseteq \text{Th}^+(\mathfrak{B}_1)$ . Both of these facts are easily checked in light of [9] Theorem VI.2.4. If  $\mathbf{a} \leq \mathbf{b}$  are in a Brouwer algebra  $\mathfrak{B}$ , then  $\mathfrak{B}/\mathbf{a}$  B<sup>+</sup>-embeds into  $\mathfrak{B}/\mathbf{b}$  by the identity. This implies that  $\text{Th}^+(\mathfrak{B}/\mathbf{b}) \subseteq \text{Th}^+(\mathfrak{B}/\mathbf{a})$ , and it follows that the  $\mathbf{a}$  for which  $\text{Th}(\mathfrak{B}/\mathbf{a}) \subseteq \text{JAN}$  is upward-closed in any Brouwer algebra  $\mathfrak{B}$ .

Our goal is to B<sup>+</sup>-embed a certain class of Brouwer algebras into  $\mathfrak{M}/\mathbf{a}$ . For any set  $X$ , let  $\text{Fr}(X)$  denote the free distributive lattice generated by  $X$  and let  $\mathbf{0} \oplus \text{Fr}(X)$  denote  $\text{Fr}(X)$  with a new

bottom element  $\mathbf{0}$ . The elements of  $\text{Fr}(X)$  are all of the form  $\sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$  where  $V$  and the  $U_v$  are finite sets of indices and the  $\mathbf{x}_u^v$  are all in  $X$  (see for example Balbes and Dwinger [1] Section V.3). For such representations,  $\sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v \leq \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$  if and only if

$$(\forall v \in V) (\exists j \in J) (\forall i \in I_j) (\exists u \in U_v) (\mathbf{x}_u^v \leq \mathbf{y}_i^j).$$

If  $\mathbf{a}, \mathbf{b} \in \text{Fr}(X)$  are such that  $\mathbf{a} \not\leq \mathbf{b}$ , then  $\mathbf{a} \rightarrow \mathbf{b}$  exists. To see this, let  $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$  and  $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$  be representations for  $\mathbf{a}$  and  $\mathbf{b}$ . Then check

$$\mathbf{a} \rightarrow \mathbf{b} = \sum \left\{ \prod_{i \in I_j} \mathbf{y}_i^j \mid (\forall v \in V) \left( \prod_{i \in I_j} \mathbf{y}_i^j \not\leq \prod_{u \in U_v} \mathbf{x}_u^v \right) \right\}.$$

If  $\mathbf{a} \geq \mathbf{b}$  are in  $\text{Fr}(X)$  for an infinite  $X$ , then  $\mathbf{a} \rightarrow \mathbf{b}$  fails to exist because in this case  $\text{Fr}(X)$  has no least element. We see then that  $\mathbf{a} \rightarrow \mathbf{b}$  exists for every  $\mathbf{a}, \mathbf{b} \in \mathbf{0} \oplus \text{Fr}(X)$ . If  $X$  is finite, then so are  $\text{Fr}(X)$  and  $\mathbf{0} \oplus \text{Fr}(X)$ . Hence both are Brouwer algebras. Let  $\text{Fr}(n)$  denote the free distributive lattice with  $n$  generators. The logic  $\text{LM} = \bigcap_{n \in \omega} \text{Th}(\mathbf{0} \oplus \text{Fr}(n))$  is called the *Medvedev logic of finite problems*. (LM is usually defined in terms of Brouwer algebras isomorphic to the  $\mathbf{0} \oplus \text{Fr}(n)$ . See [16] for details.) We take advantage of the fact that  $\text{LM}^+ = \text{IPC}^+$  [8].

If  $X$  is infinite, then  $\mathbf{0} \oplus \text{Fr}(X)$  fails to be a Brouwer algebra only because it lacks a top element. Therefore the notion of a  $B^+$ -embedding makes sense when we allow  $\mathfrak{B}_1$  or  $\mathfrak{B}_2$  to be  $\mathbf{0} \oplus \text{Fr}(X)$ . If we let  $\mathbf{0} \oplus \text{Fr}(X) \oplus \mathbf{1}$  denote  $\text{Fr}(X)$  with a new bottom element  $\mathbf{0}$  and a new top element  $\mathbf{1}$ , then  $\mathbf{0} \oplus \text{Fr}(X) \oplus \mathbf{1}$  is always a Brouwer algebra.

For any partial order  $(P, \leq_P)$ , let  $\text{Fr}(P, \leq_P)$  denote the free distributive lattice generated by  $(P, \leq_P)$ .  $\text{Fr}(P, \leq_P)$  is the quotient  $\text{Fr}(P) / \equiv$  where, for  $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} \mathbf{x}_u^v$  and  $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} \mathbf{y}_i^j$  in  $\text{Fr}(P)$ ,  $\mathbf{a} \equiv \mathbf{b}$  if and only if  $(\mathbf{a} \leq \mathbf{b}) \wedge (\mathbf{b} \leq \mathbf{a})$  and  $\mathbf{a} \leq \mathbf{b}$  if and only if

$$(\forall v \in V) (\exists j \in J) (\forall i \in I_j) (\exists u \in U_v) (\mathbf{x}_u^v \leq_P \mathbf{y}_i^j).$$

$\text{Fr}(P, \leq_P)$  is always a distributive lattice, and  $\mathbf{0} \oplus \text{Fr}(P, \leq_P) \oplus \mathbf{1}$  is always a Brouwer algebra; also see [13].

The following lemmas facilitate our embeddings. Lemma 4.3 is a slight generalization of the claim in the proof of [13] Lemma 2.3 and of [10] Lemma 6. The embedding is done in Theorem 4.4 which is nearly identical to [14] Theorem 2.11. Part of the reason for reproducing the proof here is to (hopefully) correct the notational inconsistencies in the proof in [14]. We restate [14] Theorem 2.11 for reference.

**Theorem 4.1** ([14] Theorem 2.11). *Let  $\mathbf{d} = \prod_{i=0}^n \mathbf{d}_i$  where  $\mathbf{d}_i >_{\mathbf{M}} \mathbf{0}'$  and  $\mathbf{d}_i$  is join-irreducible for each  $i \leq n$ . Then  $\mathbf{0} \oplus \text{Fr}(P, \leq_P) \oplus \mathbf{1}$  B-embeds into  $\mathfrak{M} / \mathbf{d}$  for every countable partial order  $(P, \leq_P)$ .*

(The above theorem is stated more generally in [14]. Each degree  $\mathbf{d}_i$  for  $i \leq n$  is allowed to be either join-irreducible or  $\mathfrak{D}\mathfrak{e}$ -irreducible. A degree  $\mathbf{a}$  is *dense* if it is of the form  $[\mathcal{A}]$  where  $\mathcal{A}$  is dense in  $\omega^\omega$ , and a degree  $\mathbf{d}$  is  *$\mathfrak{D}\mathfrak{e}$ -irreducible* if  $\mathbf{a} \rightarrow \mathbf{d} = \mathbf{d}$  for all dense degrees  $\mathbf{a}$ . We do not consider  $\mathfrak{D}\mathfrak{e}$ -irreducible degrees in our version of [14] Theorem 2.11, which is Theorem 4.4 below, because in Theorem 4.4 we require that the mass problems  $\mathcal{D}_i$  (which play the role of the degrees  $\mathbf{d}_i$  in [14] Theorem 2.11) are Turing upward-closed. Mass problems that are Turing upward-closed are dense and hence their degrees are not  $\mathfrak{D}\mathfrak{e}$ -irreducible.)

**Lemma 4.2** ([3]). *If  $\mathcal{X} \not\leq_{\mathbf{M}} \mathcal{Y}$  are mass problems, then there is a  $\mathcal{W} \subseteq \mathcal{X}$  with  $|\mathcal{W}| \leq \omega$  such that  $\mathcal{W} \not\leq_{\mathbf{M}} \mathcal{Y}$ .*

*Proof.*  $\mathcal{X} \not\leq_{\mathbf{M}} \mathcal{Y}$  means that there is no Turing functional  $\Phi$  such that  $\Phi(\mathcal{X}) \subseteq \mathcal{Y}$ . Thus for each functional  $\Phi_e$  there must be some function  $f_e \in \mathcal{X}$  such that  $\Phi_e(f_e) \notin \mathcal{Y}$ . Let  $\mathcal{W}$  consist of a choice of one such  $f_e \in \mathcal{X}$  for each functional  $\Phi_e$ .  $\square$



**Lemma 4.3.** *Let  $\mathcal{U}$ ,  $\mathcal{V}$ , and  $\mathcal{F}_i$  for  $i \in \omega$  be mass problems such that  $\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{F}_i \leq_M \mathcal{U} + \mathcal{V}$  and  $\sigma \hat{\wedge} \mathcal{U} \subseteq \mathcal{U}$  for all  $\sigma \in \omega^{<\omega}$ . Then there are mass problems  $\mathcal{V}_i$  for  $i \in \omega$  such that  $\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{V}_i \equiv_M \mathcal{V}$  and  $\mathcal{F}_i \leq_M \mathcal{U} + \mathcal{V}_i$  for each  $i \in \omega$ .*

*Proof.* Let  $\mathcal{U}$ ,  $\mathcal{V}$ , and  $\mathcal{F}_i$  for  $i \in \omega$  be as in the statement of the lemma. Let  $\Phi$  be such that  $\Phi(\mathcal{U} + \mathcal{V}) \subseteq \bigcup_{i \in \omega} i \hat{\wedge} \mathcal{F}_i$ . For each  $i \in \omega$ , define  $\mathcal{V}_i = \{g \in \mathcal{V} \mid (\exists \sigma \in \omega^{<\omega})(\Phi(\sigma \oplus g)(0) = i)\}$ .  $\mathcal{V} \leq_M \bigcup_{i \in \omega} i \hat{\wedge} \mathcal{V}_i$  is clear.  $\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{V}_i \leq_M \mathcal{V}$  by the reduction which, given  $g$ , searches for a  $\sigma \in \omega^{<\omega}$  such that  $\Phi(\sigma \oplus g)(0) \downarrow$  and outputs  $\Phi(\sigma \oplus g)(0) \hat{\wedge} g$ . To see  $i \hat{\wedge} \mathcal{F}_i \leq_M \mathcal{U} + \mathcal{V}_i$ , consider the reduction which, given  $f \oplus g$ , searches for a  $\sigma \in \omega^{<\omega}$  such that  $\Phi(\sigma \oplus g)(0) = i$  and outputs  $\Phi((\sigma \hat{\wedge} f) \oplus g)$ . If  $f \oplus g \in \mathcal{U} + \mathcal{V}_i$ , then such a  $\sigma$  is found,  $\sigma \hat{\wedge} f$  is in  $\mathcal{U}$ , and  $\Phi((\sigma \hat{\wedge} f) \oplus g)$  is in  $i \hat{\wedge} \mathcal{F}_i$ .  $\square$

**Theorem 4.4.** *Let  $\mathbf{d} = [\bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i]$  where  $[\mathcal{D}_i] >_M \mathbf{0}'$ ,  $[\mathcal{D}_i]$  is join-irreducible, and  $\mathcal{D}_i$  is Turing upward-closed for each  $i \in \omega$ . Then  $\mathbf{0} \oplus \text{Fr}(2^\omega) B^+$ -embeds into  $\mathfrak{M}/\mathbf{d}$ .*

*Proof.* Let  $\mathcal{D}_i$  for  $i \in \omega$  be as in the statement of the theorem, let  $\mathcal{D} = \bigcup_{i \in \omega} i \hat{\wedge} \mathcal{D}_i$ , and let  $\mathbf{d} = [\mathcal{D}]$ . Lemma 3.7 proves that  $\mathbf{d} >_M \mathbf{0}'$ . By Lemma 4.2, let  $\mathcal{A} \subseteq \{f \mid f >_T 0\}$  be a countable mass problem such that  $\mathcal{A} \not\leq_M \mathcal{D}$ . Let  $\{f_x \mid x \in 2^\omega\}$  be a collection of functions such that  $f_x \upharpoonright_T f_y$  for all  $x, y \in 2^\omega$  with  $x \neq y$  and that  $f \not\leq_T f_x$  for all  $f \in \mathcal{A}$  and  $x \in 2^\omega$ . Such a sequence can be constructed via standard recursion-theoretic techniques: build a perfect tree whose paths are Turing incomparable and do not compute any functions in  $\mathcal{A}$ . See for example [6] Section II.4. Notice that  $\mathcal{B}_{f_x} \leq_M \mathcal{A}$  (because  $\mathcal{A} \subseteq \mathcal{B}_{f_x}$ ) for each  $x \in 2^\omega$ .

Define  $G: \mathbf{0} \oplus \text{Fr}(2^\omega) \rightarrow \mathfrak{M}$  as follows. Let  $G(\mathbf{0}) = \mathbf{0}$  and let  $G(x) = [\mathcal{B}_{f_x}]$  on the generators  $x \in 2^\omega$  of  $\text{Fr}(2^\omega)$ . Then extend  $G$  to all of  $\mathbf{0} \oplus \text{Fr}(2^\omega)$  so that  $G(\sum_{v \in V} \prod_{u \in U_v} x_u^v) = \sum_{v \in V} \prod_{u \in U_v} G(x_u^v)$ .  $G$  preserves  $\mathbf{0}$ ,  $+$ , and  $\times$  by definition, and  $G$  is injective and preserves  $\rightarrow$  by Lemma 3.2 items (iii) and (iv). Hence  $G$  is a  $B^+$ -embedding (this is essentially [13] Corollary 2.5). Now define  $H: \mathbf{0} \oplus \text{Fr}(2^\omega) \rightarrow \mathfrak{M}/\mathbf{d}$  by  $H(\mathbf{a}) = G(\mathbf{a}) \times \mathbf{d}$  for all  $\mathbf{a} \in \mathbf{0} \oplus \text{Fr}(2^\omega)$ . This  $H$  is the desired  $B^+$ -embedding. By definition,  $H$  preserves  $\mathbf{0}$ ,  $+$ , and  $\times$ . We must show that  $H$  is injective and that  $H$  preserves  $\rightarrow$ .

Clearly  $H(\mathbf{a}) = \mathbf{0}$  if and only if  $\mathbf{a} = \mathbf{0}$ , so to show that  $H$  is injective we let  $\mathbf{a}, \mathbf{b} \in \text{Fr}(2^\omega)$  be such that  $H(\mathbf{a}) \leq_M H(\mathbf{b})$  and show that  $\mathbf{a} \leq \mathbf{b}$ . Let  $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} x_u^v$  be a representation for  $\mathbf{a}$  and let  $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} y_i^j$  be a representation for  $\mathbf{b}$ .  $H(\mathbf{a}) \leq_M H(\mathbf{b})$  means that

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \leq_M \sum_{j \in J} \prod_{i \in I_j} G(y_i^j) \times \mathbf{d}.$$

Therefore

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \leq_M \sum_{j \in J} \prod_{i \in I_j} G(y_i^j) = \prod \left\{ \sum_{j \in J} G(y_{\alpha(j)}^j) \mid \alpha \in \prod_{j \in J} I_j \right\}$$

where the equality is by distributivity ( $\prod_{j \in J} I_j$  denotes the Cartesian product of the  $I_j$ 's). Hence

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \leq_M \sum_{j \in J} G(y_{\alpha(j)}^j) \text{ for each } \alpha \in \prod_{j \in J} I_j.$$

Each  $\sum_{j \in J} G(y_{\alpha(j)}^j)$  is meet-irreducible by Lemma 3.2 item (ii). Also,  $\mathbf{d} \not\leq_M \sum_{j \in J} G(y_{\alpha(j)}^j)$  for each  $\alpha \in \prod_{j \in J} I_j$  because  $\sum_{j \in J} G(y_{\alpha(j)}^j) \leq_M [\mathcal{A}]$  but  $\mathbf{d} \not\leq_M [\mathcal{A}]$ . Thus

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \leq_M \sum_{j \in J} G(y_{\alpha(j)}^j) \text{ for each } \alpha \in \prod_{j \in J} I_j,$$

and this implies that

$$\sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \leq_M \prod_{j \in J} \left\{ \sum_{i \in I_j} G(y_{\alpha(j)}^j) \mid \alpha \in \prod_{j \in J} I_j \right\}.$$

The left-hand side of the above inequality is  $G(\mathbf{a})$  and the right-hand side is  $G(\mathbf{b})$ .  $G$  is a  $B^+$ -embedding, so we conclude  $\mathbf{a} \leq \mathbf{b}$ .

If either of  $\mathbf{a}, \mathbf{b} \in \mathbf{0} \oplus \text{Fr}(2^\omega)$  is  $\mathbf{0}$ , then clearly  $H(\mathbf{a} \rightarrow \mathbf{b}) = H(\mathbf{a}) \rightarrow H(\mathbf{b})$ . So as before, let  $\mathbf{a} = \sum_{v \in V} \prod_{u \in U_v} x_u^v$  and let  $\mathbf{b} = \sum_{j \in J} \prod_{i \in I_j} y_i^j$  be in  $\text{Fr}(2^\omega)$ . We see  $H(\mathbf{a} \rightarrow \mathbf{b}) \geq_M H(\mathbf{a}) \rightarrow H(\mathbf{b})$  because

$$H(\mathbf{a} \rightarrow \mathbf{b}) + H(\mathbf{a}) = H((\mathbf{a} \rightarrow \mathbf{b}) + \mathbf{a}) \geq_M H(\mathbf{b}).$$

To show that  $H(\mathbf{a} \rightarrow \mathbf{b}) \leq_M H(\mathbf{a}) \rightarrow H(\mathbf{b})$ , we show that if  $\mathbf{z} \in \mathfrak{M}$  is such that  $H(\mathbf{b}) \leq_M H(\mathbf{a}) + \mathbf{z}$ , then  $H(\mathbf{a} \rightarrow \mathbf{b}) \leq_M \mathbf{z}$ . Suppose  $H(\mathbf{b}) \leq_M H(\mathbf{a}) + \mathbf{z}$ . That is,

$$(1) \quad \sum_{j \in J} \prod_{i \in I_j} G(y_i^j) \times \mathbf{d} \leq_M \left( \sum_{v \in V} \prod_{u \in U_v} G(x_u^v) \times \mathbf{d} \right) + \mathbf{z}.$$

Since  $\mathbf{a} \rightarrow \mathbf{b} = \sum \{ \prod_{i \in I_j} y_i^j \mid (\forall v \in V) (\prod_{i \in I_j} y_i^j \not\leq_M \prod_{u \in U_v} x_u^v) \}$ , we have

$$\begin{aligned} H(\mathbf{a} \rightarrow \mathbf{b}) &= G(\mathbf{a} \rightarrow \mathbf{b}) \times \mathbf{d} \\ &= \sum \left\{ \prod_{i \in I_j} G(y_i^j) \mid (\forall v \in V) \left( \prod_{i \in I_j} G(y_i^j) \not\leq_M \prod_{u \in U_v} G(x_u^v) \right) \right\} \times \mathbf{d}. \end{aligned}$$

It suffices to show that, given  $j \in J$ , if  $\prod_{i \in I_j} G(y_i^j)$  satisfies

$$(\forall v \in V) \left( \prod_{i \in I_j} G(y_i^j) \not\leq_M \prod_{u \in U_v} G(x_u^v) \right),$$

then  $\prod_{i \in I_j} G(y_i^j) \times \mathbf{d} \leq_M \mathbf{z}$ . Suppose  $\prod_{i \in I_j} G(y_i^j)$  is such a meet. Then we know

$$(\forall v \in V) \left( \exists u \in U_v \right) \left( \prod_{i \in I_j} G(y_i^j) \not\leq_M G(x_u^v) \right).$$

By choosing such a  $u \in U_v$  for every  $v \in V$  and by appealing to Lemma 3.2 items (i) and (ii), we see that there is an  $\alpha \in \prod_{v \in V} U_v$  such that

$$(2) \quad \prod_{i \in I_j} G(y_i^j) \not\leq_M \sum_{v \in V} G(x_{\alpha(v)}^v).$$

Distributing  $\sum_{v \in V} \prod_{u \in U_v} G(x_u^v)$  in the right-hand side of (1) yields

$$\prod_{i \in I_j} G(y_i^j) \times \mathbf{d} \leq_M \sum_{v \in V} G(x_{\alpha(v)}^v) + \mathbf{z}.$$

The degree  $\sum_{v \in V} G(x_{\alpha(v)}^v)$  is a finite join of degrees of the form  $[\mathcal{B}_f]$  and thus has a representative  $\mathcal{U}$  such that  $\sigma \wedge \mathcal{U} \subseteq \mathcal{U}$  for all  $\sigma \in \omega^{<\omega}$ . So by Lemma 4.3 there are mass problems  $\mathcal{Z}_i$  for  $i \in I_j$  and

$\widehat{\mathcal{Z}}_i$  for  $i \in \omega$  such that

$$\mathbf{z} = \left( \prod_{i \in I_j} [\mathcal{Z}_i] \right) \times \left[ \bigcup_{i \in \omega} i \widehat{\mathcal{Z}}_i \right],$$

$$G(y_i^j) \leq_M \sum_{v \in V} G(x_{\alpha(v)}^v) + [\mathcal{Z}_i] \text{ for each } i \in I_j, \text{ and}$$

$$[\mathcal{D}_i] \leq_M \sum_{v \in V} G(x_{\alpha(v)}^v) + [\widehat{\mathcal{Z}}_i] \text{ for each } i \in \omega.$$

Each  $G(y_i^j)$  is join-irreducible, and  $G(y_i^j) \not\leq_M \sum_{v \in V} G(x_{\alpha(v)}^v)$  by (2). Thus  $G(y_i^j) \leq_M [\mathcal{Z}_i]$  for each  $i \in \omega$ , so  $\prod_{i \in I_j} G(y_i^j) \leq_M \prod_{i \in I_j} [\mathcal{Z}_i]$ . Each  $[\mathcal{D}_i]$  is join-irreducible by assumption, and also  $[\mathcal{D}_i] \not\leq_M \sum_{v \in V} G(x_{\alpha(v)}^v)$  because the right-hand side is  $\leq_M [\mathcal{A}]$  but the left-hand side is not. It follows that  $[\mathcal{D}_i] \leq_M [\widehat{\mathcal{Z}}_i]$  for each  $i \in \omega$ , and so  $\widehat{\mathcal{Z}}_i \subseteq \mathcal{D}_i$  for each  $i \in \omega$  because each  $\mathcal{D}_i$  is Turing upward-closed. Thus  $\bigcup_{i \in \omega} i \widehat{\mathcal{Z}}_i \subseteq \mathcal{D}$ , so  $\mathbf{d} \leq_M [\bigcup_{i \in \omega} i \widehat{\mathcal{Z}}_i]$ . Therefore

$$\prod_{i \in I_j} G(y_i^j) \times \mathbf{d} \leq_M \left( \prod_{i \in I_j} [\mathcal{Z}_i] \right) \times \left[ \bigcup_{i \in \omega} i \widehat{\mathcal{Z}}_i \right] = \mathbf{z}$$

as desired.  $\square$

**Corollary 4.5.** *If  $\mathbf{a} \geq_M \mathbf{d}$  are degrees such that  $\mathbf{d} = [\bigcup_{i \in \omega} i \mathcal{D}_i]$  where  $[\mathcal{D}_i] >_M \mathbf{0}'$  and is join-irreducible for each  $i \in \omega$ , then  $\mathbf{0} \oplus \text{Fr}(2^\omega) B^+$ -embeds into  $\mathfrak{M}/\mathbf{a}$ .*

*Proof.* Let  $\mathbf{a}$ ,  $\mathbf{d}$ , and  $\mathcal{D}_i$  for  $i \in \omega$  be as in the statement of the corollary. Let  $\mathbf{d}_0 = [\bigcup_{i \in \omega} i C(\mathcal{D}_i)]$  and notice that  $\mathbf{d} \geq_M \mathbf{d}_0$ .  $\mathcal{D}_i \equiv_M C(\mathcal{D}_i)$  for each  $i \in \omega$  by Lemma 2.3, so  $\mathbf{d}_0$  satisfies the hypotheses of Theorem 4.4. Thus  $\mathbf{0} \oplus \text{Fr}(2^\omega) B^+$ -embeds into  $\mathfrak{M}/\mathbf{d}_0$  which  $B^+$ -embeds into  $\mathfrak{M}/\mathbf{a}$ .  $\square$

**Corollary 4.6.** *If  $\mathbf{a} \geq_M \mathbf{d}$  are degrees such that  $\mathbf{d} = [\bigcup_{i \in \omega} i \mathcal{D}_i]$  where  $[\mathcal{D}_i] >_M \mathbf{0}'$  and is join-irreducible for each  $i \in \omega$ , then  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$ .*

*Proof.* The Brouwer algebra  $\mathbf{0} \oplus \text{Fr}(n) B^+$ -embeds into  $\mathbf{0} \oplus \text{Fr}(2^\omega)$  for each  $n$ , and  $\mathbf{0} \oplus \text{Fr}(2^\omega) B^+$ -embeds into  $\mathfrak{M}/\mathbf{a}$  by Corollary 4.5. Thus  $\text{Th}^+(\mathfrak{M}/\mathbf{a}) \subseteq \bigcap_{n \in \omega} \text{Th}^+(\mathbf{0} \oplus \text{Fr}(n)) = \text{LM}^+ = \text{IPC}^+$ . So  $\text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}$ .  $\square$

Theorem 4.4 can be modified to  $B$ -embed  $\mathbf{0} \oplus \text{Fr}(2^\omega) \oplus \mathbf{1}$  into  $\mathfrak{M}/\mathbf{d}$  for degrees  $\mathbf{d}$  as in the statement of Theorem 4.4. However, if  $\mathbf{a} \leq \mathbf{b}$  in a Brouwer algebra  $\mathfrak{B}$ , it is not in general the case that  $\mathfrak{B}/\mathbf{a}$   $B$ -embeds into  $\mathfrak{B}/\mathbf{b}$ . So the proof of Corollary 4.5 fails for  $B$ -embedding  $\mathbf{0} \oplus \text{Fr}(2^\omega) \oplus \mathbf{1}$ . Theorem 4.4 can also be modified to prove a more precise analogue of [14] Theorem 2.11 (restated as Theorem 4.1 above). Let  $\mathbf{d} = [\bigcup_{i \in \omega} i \mathcal{D}_i]$  where  $[\mathcal{D}_i] >_M \mathbf{0}'$ ,  $[\mathcal{D}_i]$  is join-irreducible, and  $\mathcal{D}_i$  is Turing upward-closed for each  $i \in \omega$ . Then  $\mathbf{0} \oplus \text{Fr}(P, \leq_P) \oplus \mathbf{1}$   $B$ -embeds into  $\mathfrak{M}/\mathbf{d}$  for every countable partial order  $(P, \leq_P)$ .

## 5. $\mathfrak{F}_{\text{cl}}$ IS NOT PRIME

Recall that a filter  $\mathfrak{F}$  in a lattice is called *prime* if  $\mathbf{a} + \mathbf{b} \in \mathfrak{F} \rightarrow \mathbf{a} \in \mathfrak{F} \vee \mathbf{b} \in \mathfrak{F}$  for all  $\mathbf{a}$  and  $\mathbf{b}$  in the lattice. Theorem 2.4 can be phrased as a characterization of the prime principal filters in  $\mathfrak{M}$ : a degree  $\mathbf{a}$  generates a prime filter if and only if  $\mathbf{a} = [\omega^\omega - \mathcal{I}]$  for some Turing ideal  $\mathcal{I}$ . There is little general theory of the non-principal filters in  $\mathfrak{M}$ , but several specific cases have been studied in Dymant [3], Sorbi [11], Bianchini and Sorbi [2], and Lewis, Shore, and Sorbi [7]. See also [15] for a summary of many of the results appearing in these works. We consider the filters  $\mathfrak{F}$  and  $\mathfrak{F}_{\text{cl}}$ :

**Definition 5.1.**

- A degree  $\mathbf{a}$  is called *dense* (*closed*) if  $\mathbf{a} = [\mathcal{A}]$  for an  $\mathcal{A}$  that is dense (closed) in  $\omega^\omega$ .
- $\mathfrak{J}$  denotes the ideal generated by  $\{\mathbf{a} \mid \mathbf{a} \text{ is dense}\}$ .

- $\mathfrak{F}$  denotes  $\mathfrak{M} - \mathfrak{I}$ .
- $\mathfrak{F}_{\text{cl}}$  denotes the filter generated by  $\{\mathbf{a} \mid \mathbf{a} \succ_{\mathfrak{M}} \mathbf{0} \text{ and is closed}\}$ .

The join and meet of two dense degrees is dense [3], and the join and meet of two closed degrees is closed [2]. Thus  $\mathfrak{I} = \{\mathbf{b} \mid (\exists \mathbf{a} \geq_{\mathfrak{M}} \mathbf{b})(\mathbf{a} \text{ is dense})\}$  and  $\mathfrak{F}_{\text{cl}} = \{\mathbf{b} \mid (\exists \mathbf{a} \leq_{\mathfrak{M}} \mathbf{b})(\mathbf{a} \succ_{\mathfrak{M}} \mathbf{0} \text{ and is closed})\}$ . The basic properties of  $\mathfrak{I}$ ,  $\mathfrak{F}$ , and  $\mathfrak{F}_{\text{cl}}$  are as follows:  $\mathfrak{I}$  is a prime ideal [11],  $\mathfrak{F}$  is a prime filter [2],  $\mathfrak{I}$  is not principal [3],  $\mathfrak{F}$  and  $\mathfrak{F}_{\text{cl}}$  are not principal [2], and  $\mathfrak{F}_{\text{cl}} \subsetneq \mathfrak{F}$  [2]. Both [2] and [15] ask for a proof that  $\mathfrak{F}_{\text{cl}}$  is not prime. We provide a proof of this fact now.

**Lemma 5.2.** *For any  $f \in \omega^\omega$  there are  $\mathcal{A}, \mathcal{B} \subseteq \omega^\omega$  such that  $\mathcal{A} + \mathcal{B} \geq_{\mathfrak{M}} \{f\}$  and, for any closed  $\mathcal{C} \subseteq \omega^\omega$ , if  $\mathcal{A} \geq_{\mathfrak{M}} \mathcal{C}$  or  $\mathcal{B} \geq_{\mathfrak{M}} \mathcal{C}$ , then  $\mathcal{C}$  contains a recursive function.*

*Proof.* Fix a recursive bijection  $\omega \leftrightarrow \omega^{<\omega}$ . For  $e, n \in \omega$ , if

$$\forall m \forall \sigma (\exists \tau \supseteq \sigma) (\Phi_e(n \hat{\ } \tau)(m) \downarrow),$$

then define  $\eta(e, n, i) \in \omega^{<\omega}$  by induction on  $i \in \omega$  as follows. Let  $\eta(e, n, 0) = n \hat{\ } \sigma$ , where  $\sigma$  is the least string such that  $\Phi_e(n \hat{\ } \sigma)(0) \downarrow$ . Given  $\eta(e, n, i)$ , let  $\eta(e, n, i + 1) = \eta(e, n, i) \hat{\ } 0 \hat{\ } \sigma$ , where  $\sigma$  is the least string such that  $\Phi_e(\eta(e, n, i) \hat{\ } 0 \hat{\ } \sigma)(i + 1) \downarrow$ .

Let  $f \in \omega^\omega$ . We construct  $\mathcal{A}$  and  $\mathcal{B}$  such that:

- If  $g \in \mathcal{A}$ , then  $g(0)$  has the form

$$g(0) = \langle \ell, \langle n_0, x_0, y_0 \rangle, \dots, \langle n_{\ell-1}, x_{\ell-1}, y_{\ell-1} \rangle \rangle,$$

where  $\ell \in \omega$  and  $n_i \in \omega$ ,  $x_i \in \{0, 1\}$ , and  $y_i \in \omega$  for each  $i < \ell$ .

- If  $g \in \mathcal{A}$  and  $\langle n_e, 0, y_e \rangle$  is in the  $e^{\text{th}}$  position of  $g(0)$ , then
  - $\exists m \exists \sigma (\forall \tau \supseteq \sigma) (\Phi_e(n_e \hat{\ } \tau)(m) \uparrow)$
  - Any  $h \in \mathcal{B}$  with  $h(0) = n_e$  is of the form  $h = n_e \hat{\ } \sigma \hat{\ } f$ , where  $|\sigma| = y_e$ .
- If  $g \in \mathcal{A}$  and  $\langle n_e, 1, y_e \rangle$  is in the  $e^{\text{th}}$  position of  $g(0)$ , then
  - $\forall m \forall \sigma (\exists \tau \supseteq \sigma) (\Phi_e(n_e \hat{\ } \tau)(m) \downarrow)$
  - Any  $h \in \mathcal{B}$  with  $h(0) = n_e$  is of the form  $h = \eta(e, n_e, i) \hat{\ } 1 \hat{\ } f$  for some  $i \in \omega$ .
- The above properties hold with the roles of  $\mathcal{A}$  and  $\mathcal{B}$  reversed.

We construct  $\mathcal{A}$  and  $\mathcal{B}$  in stages. The construction is similar to the construction in Lemma 2.3 in that if  $g$  goes into  $\mathcal{A}$  before  $h$  goes into  $\mathcal{B}$ , then  $h(0)$  codes how to recover  $f$  from  $g$ , and similarly with the roles of  $\mathcal{A}$  and  $\mathcal{B}$  reversed. Start at stage 0 with  $\mathcal{A} = \emptyset$ ,  $\mathcal{B} = \emptyset$ ,  $s = \langle \rangle$ , and  $t = \langle \rangle$ .

Stage  $e + 1$ : Set  $n_e = e \hat{\ } t$ .

Case 1:  $\exists m \exists \sigma (\forall \tau \supseteq \sigma) (\Phi_e(n_e \hat{\ } \tau)(m) \uparrow)$ . Choose such a  $\sigma$  and put  $n_e \hat{\ } \sigma \hat{\ } f$  in  $\mathcal{A}$ . Update  $s = s \hat{\ } \langle n_e, 0, |\sigma| \rangle$ .

Case 2:  $\forall m \forall \sigma (\exists \tau \supseteq \sigma) (\Phi_e(n_e \hat{\ } \tau)(m) \downarrow)$ . Put the functions  $\eta(e, n_e, i) \hat{\ } 1 \hat{\ } f$  in  $\mathcal{A}$  for each  $i \in \omega$ . Update  $s = s \hat{\ } \langle n_e, 1, 0 \rangle$ .

Repeat the above procedure with the roles of  $\mathcal{A}$  and  $\mathcal{B}$  reversed and the roles of  $s$  and  $t$  reversed. This completes stage  $e + 1$ . Then go on to stage  $e + 2$ . This completes the construction.

Suppose  $\mathcal{A} \geq_{\mathfrak{M}} \mathcal{C}$  where  $\mathcal{C}$  is closed. We show that  $\mathcal{C}$  contains a recursive function. The proof with  $\mathcal{B}$  in place of  $\mathcal{A}$  is the same. Let  $\Phi_e(\mathcal{A}) \subseteq \mathcal{C}$ . Consider stage  $e + 1$  of the above construction. Case 1 must not have occurred because otherwise  $\mathcal{A}$  would contain a function  $n_e \hat{\ } \sigma \hat{\ } f$  such that  $\Phi_e(n_e \hat{\ } \sigma \hat{\ } f)$  is not total. Thus case 2 occurred, and so  $\mathcal{A}$  contains the function  $\eta(e, n_e, i) \hat{\ } 1 \hat{\ } f$  for each  $i \in \omega$ . Let  $k$  be the recursive function  $k = n_e \hat{\ } \sigma_0 \hat{\ } 0 \hat{\ } \sigma_1 \hat{\ } 0 \hat{\ } \sigma_2 \hat{\ } 0 \hat{\ } \dots$ , where  $\eta(e, n_e, i) = n_e \hat{\ } \sigma_0 \hat{\ } 0 \hat{\ } \dots \hat{\ } 0 \hat{\ } \sigma_i$  for each  $i \in \omega$  (think of  $k$  as the “limit” of the strings  $\eta(e, n_e, i)$  as  $i \rightarrow \infty$ ). Then  $\Phi_e(\eta(e, n_e, i) \hat{\ } 1 \hat{\ } f) \in \mathcal{C}$  and  $\Phi_e(\eta(e, n_e, i) \hat{\ } 1 \hat{\ } f) \upharpoonright i = \Phi_e(k) \upharpoonright i$  for each  $i \in \omega$ . Thus  $\mathcal{C}$  contains the recursive function  $\Phi_e(k)$  because  $\mathcal{C}$  is closed.

We now describe a uniform procedure for producing  $f$  from  $g \oplus h \in \mathcal{A} + \mathcal{B}$ . First decode  $h(0)$  as  $h(0) = \langle \ell, \langle n_0, x_0, y_0 \rangle, \dots, \langle n_{\ell-1}, x_{\ell-1}, y_{\ell-1} \rangle \rangle$  and look for  $g(0)$  among the  $n_e$ . If  $\langle g(0), 0, y_e \rangle$  appears in  $h(0)$  at position  $e$ , then output  $g$  from position  $y_e + 1$  onward as in this case  $g = \sigma \hat{\ } f$  for

some string  $\sigma$  of length  $y_e + 1$ . If  $\langle g(0), 1, 0 \rangle$  appears in  $h(0)$  at position  $e$ , then  $g = \eta(e, g(0), i) \wedge 1 \wedge f$  for some  $i \in \omega$ . Compute which  $i$  by successively computing the  $\eta(e, g(0), j)$ , matching them against  $g$ , and checking if the next bit of  $g$  is 0 (in which case compute  $\eta(e, g(0), j + 1)$ ) or 1 (in which case  $j = i$ ). Output  $f$  once  $i$  is found.

The number  $g(0)$  appears among the  $n_e$  coded into  $h(0)$  if  $g$  went into  $\mathcal{A}$  before  $h$  went into  $\mathcal{B}$ . Otherwise  $h$  went into  $\mathcal{B}$  before  $g$  went into  $\mathcal{A}$ , so  $h(0)$  appears among the  $n_e$  coded in  $g(0)$ . In this case, switch the roles of  $g$  and  $h$  and apply the above procedure to compute  $f$ .  $\square$

**Theorem 5.3.**  $\mathfrak{F}_{\text{cl}}$  is not prime. In fact, if  $\mathfrak{G} \subseteq \mathfrak{F}_{\text{cl}}$ ,  $\mathfrak{G} \neq \{\mathbf{1}\}$  is a filter, then  $\mathfrak{G}$  is not prime.

*Proof.* Suppose  $\mathfrak{G} \subseteq \mathfrak{F}_{\text{cl}}$  is a filter such that  $\mathfrak{G} \neq \{\mathbf{1}\}$ . Let  $f >_{\text{T}} \mathbf{0}$  be such that  $[\{f\}] \in \mathfrak{G}$ . Let  $\mathcal{A}, \mathcal{B} \subseteq \omega^\omega$  be as in Lemma 5.2 for this  $f$ . Let  $\mathbf{a} = [\mathcal{A}]$  and  $\mathbf{b} = [\mathcal{B}]$ . Then  $\mathbf{a}, \mathbf{b} \notin \mathfrak{G}$  because  $\mathbf{a}, \mathbf{b} \notin \mathfrak{F}_{\text{cl}}$ , but  $\mathbf{a} + \mathbf{b} \in \mathfrak{G}$  because  $\mathbf{a} + \mathbf{b} \geq_{\text{M}} [\{f\}]$ .  $\square$

If  $\mathbf{x}$  and  $\mathbf{y}$  are degrees such that  $\mathbf{y}$  is closed and  $\mathbf{y} \not\leq_{\text{M}} \mathbf{x}$ , then there is no dense degree  $\mathbf{z}$  such that  $\mathbf{y} \leq_{\text{M}} \mathbf{x} + \mathbf{z}$  [7]. Therefore, if  $\mathfrak{G} \subseteq \mathfrak{F}_{\text{cl}}$ ,  $\mathfrak{G} \neq \{\mathbf{1}\}$  is a filter, then any degrees  $\mathbf{a}$  and  $\mathbf{b}$  witnessing that  $\mathfrak{G}$  is not prime must both be in  $\mathfrak{F} - \mathfrak{G}$ .

The results of Section 3 suggest two new filters to study:

**Definition 5.4.**

- $\mathfrak{G}$  denotes the filter generated by

$$\{\mathbf{d} \mid \mathbf{d} >_{\text{M}} \mathbf{0}' \text{ and is join-irreducible}\}.$$

- $\mathfrak{H}$  denotes the filter generated by

$$\left\{ \left[ \bigcup_{i \in \omega} i \wedge \mathcal{D}_i \right] \mid (\forall i \in \omega) ([\mathcal{D}_i] >_{\text{M}} \mathbf{0}' \text{ and is join-irreducible}) \right\}.$$

$\mathfrak{G}$  is exactly the set of all degrees  $\mathbf{b}$  for which  $\mathbf{b} \geq_{\text{M}} \prod_{i=0}^n \mathbf{d}_i$  for some join-irreducible degrees  $\mathbf{d}_i >_{\text{M}} \mathbf{0}'$ ,  $i \leq n$ , and  $\mathfrak{H}$  is exactly the set of all degrees  $\mathbf{b}$  for which  $\mathbf{b} \geq_{\text{M}} [\bigcup_{i \in \omega} i \wedge \mathcal{D}_i]$  for some join-irreducible degrees  $[\mathcal{D}_i] >_{\text{M}} \mathbf{0}'$ ,  $i \in \omega$ .

**Theorem 5.5.**  $\mathfrak{F}_{\text{cl}} \subsetneq \mathfrak{G} \subsetneq \mathfrak{H} \subsetneq \{\mathbf{a} \mid \mathbf{a} >_{\text{M}} \mathbf{0}'\}$ .  $\mathfrak{G} \not\subseteq \mathfrak{F}$  (hence also  $\mathfrak{H} \not\subseteq \mathfrak{F}$ ). Neither  $\mathfrak{G}$  nor  $\mathfrak{H}$  is principal.

*Proof.* Every closed degree  $>_{\text{M}} \mathbf{0}$  bounds a join-irreducible degree  $>_{\text{M}} \mathbf{0}'$  [16]. Hence  $\mathfrak{F}_{\text{cl}} \subseteq \mathfrak{G}$ .  $\mathfrak{G} \subseteq \mathfrak{H}$  is clear. To see  $\mathfrak{G} \not\subseteq \mathfrak{F}$ , observe that every  $\mathcal{B}_f$  is dense, so if  $f >_{\text{T}} \mathbf{0}$ , then  $[\mathcal{B}_f] \in \mathfrak{G} - \mathfrak{F}$ . This also shows  $\mathfrak{G} \not\subseteq \mathfrak{F}_{\text{cl}}$ . The degree constructed in Theorem 3.8 witnesses  $\mathfrak{H} \not\subseteq \mathfrak{G}$ . The degree constructed in Theorem 3.10 witnesses  $\{\mathbf{a} \mid \mathbf{a} >_{\text{M}} \mathbf{0}'\} \not\subseteq \mathfrak{H}$ . We show that  $\mathfrak{G}$  is not principal. The proof for  $\mathfrak{H}$  is the same. First, if  $\mathcal{A}$  is countable and contains no recursive functions, then there is a function  $f >_{\text{T}} \mathbf{0}$  such that  $g \not\leq_{\text{T}} f$  for all  $g \in \mathcal{A}$ . Thus  $\mathcal{B}_f \leq_{\text{M}} \mathcal{A}$  (as  $\mathcal{A} \subseteq \mathcal{B}_f$ ) for this  $f$ . Every  $[\mathcal{B}_f]$  for  $f >_{\text{T}} \mathbf{0}$  is in  $\mathfrak{G}$ , so every  $[\mathcal{A}]$  where  $\mathcal{A}$  is countable and contains no recursive function is in  $\mathfrak{G}$ . If  $\mathfrak{G}$  were principal, it would be generated by a degree  $\mathbf{x}$  such that  $\mathbf{x} \leq_{\text{M}} [\mathcal{A}]$  for all countable  $\mathcal{A}$  not containing a recursive function. By Lemma 4.2, the only such  $\mathbf{x}$  are  $\mathbf{0}$  and  $\mathbf{0}'$ . We know  $\mathbf{0}$  and  $\mathbf{0}'$  are not in  $\mathfrak{G}$ , so  $\mathfrak{G}$  cannot be principal.  $\square$

We end with a question.

**Question 5.6.**

- Is  $\mathfrak{F} \subseteq \mathfrak{G}$ ? Is  $\mathfrak{F} \subseteq \mathfrak{H}$ ?
- Is  $\mathfrak{G}$  prime? Is  $\mathfrak{H}$  prime?
- Is  $\{\mathbf{a} \mid \text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}\}$  a filter?

To prove that  $\{\mathbf{a} \mid \text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}\}$  is a filter, it suffices to prove that  $\text{Th}(\mathfrak{M}/(\mathbf{a} \times \mathbf{b})) \subseteq \text{JAN}$  whenever both  $\text{Th}(\mathfrak{M}/\mathbf{a})$  and  $\text{Th}(\mathfrak{M}/\mathbf{b})$  are  $\subseteq \text{JAN}$  because  $\{\mathbf{a} \mid \text{Th}(\mathfrak{M}/\mathbf{a}) \subseteq \text{JAN}\}$  is upward-closed in  $\mathfrak{M}$ .

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